

Advanced Features of Macro Instructions

This lecture will focus on some of the advanced features of the macro language as implemented by the IBM/System 360 assembler.

We shall focus on our stack handling macros.

Some of the features to be covered by this lecture include.

1. The use of concatenation to generate type-specific instructions.
2. Some standard system variable symbols.
3. The use of one system variable symbol to solve the branch problem.
4. Conditional assembly.
5. The use of conditional assembly as a help in writing **STKPOP**.
7. The **ABEND** macro and its use in signaling run-time errors.
8. A completed version of our stack macros.

Concatenation: Building Operations

In a model statement, it is possible to concatenate two strings of characters.

Consider the macro prototype to load a register from one of several sources. Note the use of the string “&NAME” to allow this to be a branch target.

```
MACRO
&NAME    LOAD &REG , &TYPE , &ARG
&NAME    L&TYPE &REG , &ARG
MEND
```

Consider a number of invocations.

```
LOAD R7 , R , R6  becomes    LR R7 , R6
```

```
LOAD R7 , H , HW  becomes    LH R7 , HW
```

```
LOAD R7 , , FW    becomes    L R7 , FW
```

Note here: the second argument is empty. The empty string is concatenated to “F”.

We shall now extend the stack operations to push and pop contents of half-words and full-words, as well as registers.

Pushing from Various Sources

We look first at the handling of our **STKPUSH**. The only restriction on the stack is that every value pushed be treated as a 32-bit fullword.

As a result, a 16-bit halfword will be sign-extended to a 32-bit fullword before being pushed onto the stack. This is similar to the function of the **LH** instruction, which loads a register from a halfword.

The key instruction in the original **STKPUSH** macro is the following.

```
ST    &R,0(3,2)        STORE THE ITEM INTO THE STACK
```

In this case, the item to be placed on the stack is found in the register indicated by the symbolic parameter **&R**.

The way to extend this instruction to all data types is as follows.

1. Select a register to be a fixed source for the word on the stack, and
2. Construct instructions to load that fixed register from the source.

What Shall Be Stored on the Stack?

At this point, we have a decision to make. What data types to store?

The size restriction on the stack limits the simple choices to addresses and the contents of registers, halfwords, and fullwords.

We must select a working register for the new macro. I select R4.

The “key code” becomes as follows.

Stacking an address	LA R4 ,&ARG	Load address into R4.
	ST R4 ,&R,0(3,2)	

Stacking a halfword	LH R4 ,&ARG	Load halfword into R4.
	ST R4 ,&R,0(3,2)	

Stacking a fullword	L R4 ,&ARG	Load fullword into R4.
	ST R4 ,&R,0(3,2)	

Stacking a register	LR R4 ,&ARG	Load value from source register into R4.
	ST R4 ,&R,0(3,2)	

Passing the Type in a Macro Invocation

The solution adopted to the problem above is to pass the type in the macro call and use concatenation to build the load operator.

Here is some code taken from a macro definition that has been run and tested. First, we show the macro prototype.

```
&L2           STKPUSH &ARG , &TYP
```

Next we show the “key instruction” in the macro body.

```
L&TYP R4 , &ARG
```

Here are four typical invocations of the macro.

```
STKPUSH R7 , R           PUSH VALUE IN REGISTER.  
STKPUSH HHW , H         PUSH A HALFWORD VALUE.  
STKPUSH FFW , A         PUSH AN ADDRESS.  
STKPUSH FFW             PUSH A FULLWORD.
```

Note that the last invocation lacks a second argument. In the expansion, this causes **&TYP** to be set to `\ ’`, a blank; “**L&TYP**” becomes “**L ’**”.

The Macro Definition

Here is the definition for the macro at this stage of its development.

```
MACRO
&L2    STKPUSH &ARG,&TYP
&L2    LH      R3,STKCOUNT
        SLA    R3,2
        LA     R2,THESTACK
        L&TYP  R4,&ARG
        ST     R4,0(3,2)
        LH     R3,STKCOUNT
        AH     R3,=H'1'
        STH    3,STKCOUNT
MEND
```

Again, the “&L2” allows the macro invocation to be a branch target.

Some Invocations of this Macro

```
91      STKPUSH R7,R
92+    LH      R3,STKCOUNT
93+    SLA     R3,2
94+    LA      R2,THESTACK
95+    LR      R4,R7
96+    ST      R4,0(3,2)
97+    LH      R3,STKCOUNT
98+    AH      R3,=H'1'
99+    STH     3,STKCOUNT
```

```
100     STKPUSH HHW,H
101+    LH      R3,STKCOUNT
102+    SLA     R3,2
103+    LA      R2,THESTACK
104+    LH      R4,HHW
105+    ST      R4,0(3,2)
106+    LH      R3,STKCOUNT
107+    AH      R3,=H'1'
108+    STH     3,STKCOUNT
```

More Invocations of this Macro

```
109      STKPUSH FFW
110+     LH      R3,STKCOUNT
111+     SLA     R3,2
112+     LA      R2,THESTACK
113+     L       R4,FFW
114+     ST      R4,0(3,2)
115+     LH      R3,STKCOUNT
116+     AH      R3,=H'1'
117+     STH     3,STKCOUNT
```

```
118      STKPUSH FFW,A
119+     LH      R3,STKCOUNT
120+     SLA     R3,2
121+     LA      R2,THESTACK
122+     LA      R4,FFW
123+     ST      R4,0(3,2)
124+     LH      R3,STKCOUNT
125+     AH      R3,=H'1'
126+     STH     3,STKCOUNT
```

NOTE: The originals of the program listing are found at the end of the slides.

Saving the Work Registers

As written, this macro has the side effect of changing the values of three registers: R2, R3, & R4. The value of R4 is preserved only if it is being pushed.

We should write macros so that they operate without side effects. The only way to do this is to save and restore the values of the work registers.

There are many ways to do this. The simplest is to alter the stack data structure. Here is the new version.

```
STKCOUNT DC H`0'      NUMBER OF ITEMS STORED ON STACK
STKSIZE   DC H`64'     MAXIMUM STACK CAPACITY
STKSAV2   DC F`0'     SAVES CONTENTS OF R2
STKSAV3   DC F`0'     SAVES CONTENTS OF R3
STKSAV4   DC F`0'     SAVES CONTENTS OF R4
THESTACK  DC 64F`0'   THE STACK HOLDS 64 FULLWORDS
```

This new definition does not alter the **STKINIT** macro. It does affect the other two macros: **STKPOP** and **STKPUSH**. We illustrate the latter.

The First Revision of STKPUSH

Here is the revision that allows the work registers to be saved.

```
MACRO
&L2    STKPUSH  &ARG,&TYP
&L2    ST      R2,STKSAV2    THE ORDER OF SAVING
      ST      R3,STKSAV3    IS NOT IMPORTANT.
      ST      R4,STKSAV4
      LH      R3,STKCOUNT
      SLA     R3,2
      LA      R2,THESTACK
      L&TYP  R4,&ARG
      ST      R4,0(3,2)
      LH      R3,STKCOUNT
      AH      R3,=H'1'
      STH     R3,STKCOUNT
      L       R4,STKSAV4    THE ORDER OF RESTORATION
      L       R3,STKSAV3    IS NOT IMPORTANT EITHER.
      L       R2,STKSAV2
MEND
```

The Status of the Macros at This Point

There are a few issues to be addressed at this point.

The only macro that will not change is the initialization macro, **STKINIT**.

1. We have not yet dealt with generalizing the **STKPOP** macro.
2. We have not yet dealt with either the stack empty problem or that of the stack being full. Each has to be addressed.

Each of these issues demands the use of techniques we have not yet discussed.

Consider the first problem. We shall want to pop the following from the stack: register values, halfwords, and fullwords. The type for the argument refers to the destination; an address can be popped into either a register or fullword.

In order to see the problem for **STKPOP**, consider the “key instruction”.

Halfword: **STH R4 , &ARG**

Fullword: **ST R4 , &ARG**

Register: **LR &ARG , R4 No STR for store register.**

We could write a **STR** macro, but I want to use another solution.

Some System Variable Symbols

There are a number of system variable symbols. I mention three.

- &SYSDATE** The system date, in the 8 character form “**MM/DD/YY**”.
Use in the form of a declaration of initialized storage, as in
TODAY DC C`&SYSDATE`
- &SYSTIME** The system time of day, in the five character form “**HH.MM**”.
Also used in the form of a declaration, as in
NOW DC C`&SYSTIME`
- &SYSNDX** The macro expansion index. For the first macro expansion,
the Assembler initializes **&SYSNDX** to the string “**0001**”.
Each macro invocation increases the value represented by 1,
giving rise to the sequence “**0001**”, “**0002**”, “**0003**”, etc.

The **&SYSNDX** system variable symbol can prevent a macro from generating duplicate labels. The system symbol is concatenated to a leading character, which begins the label and must be unique within the macro definition.

More on the Macro Expansion Index

First consider the following string, used as a label in a macro definition.

L&SYSNDX L R4,STKSAV4

Note that the string “**L&SYSNDX**”, as written, contains eight characters: the initial character “**L**” followed by the 7 character sequence “**&SYSNDX**”.

On expansion, this will be converted to labels such as “**L0001**”, “**L0002**”, etc.

In the macro definition, this takes the maximum eight characters allowed for a properly formatted listing. For this reason, I suggest that the better form for the label in the macro definition is **Single_Letter&SYSNDX**.

In actual fact, the requirement for the leading characters, to which the **&SYSNDX** is to be appended can be any sequence of one to four characters, provided only that the first character is a letter. Thus the following are valid.

A12&SYSNDX ... This label might become A120003.

WXYZ&SYSNDX ... This might become WXYZ0117.

I suggest use of a single leading letter, this allows 26 labels per macro.

A Simple Example of Label Generation

Consider the simple macro used for packed division in the previous lecture. We adapt it to prevent division by zero.

```
MACRO
&LABEL DIVID &QUOT , &DIVIDEND , &DIVISOR
&LABEL ZAP &QUOT , &DIVIDEND
CP &DIVISOR , =P ` 0 ' IS IT ZERO
BNE A&SYSNDX NO , DIVISION IS OK
ZAP &QUOT , =P ` 0 ' YES , SET QUOTIENT TO 0
B B&SYSNDX
A&SYSNDX DP &QUOT , &DIVISOR
B&SYSNDX NOPR R3 DO NOTHING
MEND
```

Note that the format of the **NOPR** instruction requires a register number (here **R3**), even though the instruction does nothing.

Sample Expansion of the Macro

With the above definition, consider the following expansions.

```
A10START DIVID X,Y,Z
+A10START ZAP X,Y
+ CP Z,=P`0' IS IT ZERO
+ BNE A0001 NO, DIVISION IS OK
+ ZAP X,=P`0' YES, SET QUOTIENT TO 0
+ B B0001
+A0001 DP X,Z
+B0001 NOPR R3 DO NOTHING

A20DOIT DIVID A,B,C
+A20DOIT ZAP A,B
+ CP C,=P`0' IS IT ZERO
+ BNE A0002 NO, DIVISION IS OK
+ ZAP X,=P`0' YES, SET QUOTIENT TO 0
+ B B0002
+A0002 DP A,C
+B0002 NOPR R3 DO NOTHING
```

Note that each invocation has distinct labels. This removes the name clashes.

Another Design Strategy for DIVID

In this variant, a zero divisor will cause the program to terminate abnormally.

```
MACRO
&LABEL DIVID &QUOT , &DIVIDEND , &DIVISOR
&LABEL ZAP &QOUT , &DIVIDEND
CP &DIVISOR , =P `0' IS IT ZERO
BNE A&SYSNDX NO , DIVISION IS OK
ABEND INVOKE THE MACRO TO
TERMINATE EXECUTION.
A&SYSNDX DP &QUOT , &DIVISOR
MEND
```


The First Revision of STKINIT

Here is a revision of the STKINIT code that allows initialization of its size.

```
35          MACRO
36 &L1      STKINIT &SIZE
37 &L1      ST R3,STKSAV3
38          SR R3,R3
39          STH R3,STKCOUNT
40          L  R3,STKSAV3
41          B  L&SYSNDX
42 STKCOUNT DC H'0'
43 STKSIZE   DC H'&SIZE'
44 STKSAV2   DC F'0'
45 STKSAV3   DC F'0'
46 STKSAV4   DC F'0'
47 THESTACK DC &SIZE.F'0'
48 L&SYSNDX  SLA R3,0
49          MEND
```

Note the “.” in the definition of **THESTACK**. This concatenates the value of the symbolic parameter with “F’0’”, as in “128F’0’”

The Second Revision of STKPUSH

```
MACRO
&L2 STKPUSH &ARG,&TYP
&L2 ST R3,STKSAV3
LH R3,STKCOUNT GET COUNT OF ITEMS ON THE STACK
CH R3,STKSIZE IS THE STACK FULL?
BNL Z&SYSNDX YES, DO NOT ADD ANOTHER.
ST R4,STKSAV4 NO, WE CAN PUSH ANOTHER ITEM.
ST R2,STKSAV2 START BY SAVING THE OTHER 2 REGISTERS
SLA R3,2 MULTIPLY THE INDEX BY 4.
LA R2,THESTACK
L&TYP R4,&ARG FORM THE ADDRESS
ST R4,0(3,2) STORE THE ITEM
LH R3,STKCOUNT GET THE OLD COUNT OF ITEMS
AH R3,=H'1' INCREMENT THE COUNT BY 1
STH R3,STKCOUNT STORE THE CURRENT COUNT
L R4,STKSAV4 RESTORE THE REGISTERS.
L R2,STKSAV2
Z&SYSNDX L R3,STKSAV3
MEND
```

This is the final version of the **STKPUSH** macro.

We must discuss another basic topic before addressing **STKPOP**.

Conditional Assembly

We have already seen how concatenation can be used to construct different instructions in a macro expansion.

We now investigate conditional assembly, in which the expansion of a macro can lead to a number of distinct code sequences.

Conditional assembly permits the testing of attributes such as data format, data value, or field length, and to use the results of such testing to generate source code that is specific to the case in question.

This lecture will focus on five specific conditional assembly instructions.

AGO an unconditional branch

AIF a conditional branch. This means “Ask If”.

ANOP A NOP that can be the branch target for either **AGO** or **AIF**.

MNOTE print a programmer defined message at assembly time

MEXIT exit the macro definition.

Attributes for Use by Conditional Assembly

The assembler can generate code specified by certain attributes of the arguments to the macro definition at the time it is expanded.

There are six types of attributes that can be associated with a parameter. Here are three of the more useful attributes.

L'	Length	The length of the symbolic parameter
I'	Integer	The integer attribute of a fixed-point, floating-point, or packed decimal number.
T'	Type	The type of the parameter, as specified by the DC or DS declaration with which it is defined.

Some types for the T' attribute are as follows.

A	Address	C	Character	H	Halfword	P	Packed Decimal
B	Binary	F	Fullword	I	Instruction	X	Hexadecimal

The Ask If (AIF) Instruction

The **AIF** instruction has two parts.

1. A logical expression in parentheses, and
2. A sequence symbol immediately following, which serves as the branch target.

The **AIF** logical expression may use the following relational operators, which are quite similar to those seen in early versions of the FORTRAN language.

EQ	Equal To	NE	Not Equal To
LT	Less Than	GE	Greater Than or Equal To
GT	Greater Than	LE	Less Than or Equal To

If the type of **&AMT** is packed, go to **.B23PACK**

```
AIF(T'&AMT EQ 'P').B23PACK
```

If the type of **&LINK** is not an instruction, go to **.R30ERROR**

```
AIF(T'&LINK NE 'I').R30ERROR
```

Testing the Value of a Symbolic Parameter

What we want for the STKPOP instruction is a conditional assembly based on the value of the second parameter.

The prototype will be something like

```
&L1      STKPOP &ARG,&TYP
```

What we want to issue is an **AIF** statement such as

```
AIF (&TYP EQ 'R').ISREG
```

There is a well-known peculiarity in assembler language, not just in the IBM Assembler, that disallows this straightforward construct.

We must put the symbolic parameter in single quotes. The statement is thus:

```
AIF ('&TYP' EQ 'R').ISREG
```

If **&TYP** is the character R, the logical expression becomes ('R' EQ 'R'), which immediately evaluates to True, and the branch is taken.

Reference

Page 384, High Level Assembler for z/OS & z/VM & z/VSE Language Reference Manual, Release 6 (July 2008), SC26-4940-05

Targets for Use by Conditional Assembly

Each of the **AGO** and **AIF** instructions is a branch instruction that takes effect at assembly time. Neither persists into the assembly source code.

It should be expected that the targets for either of these conditional assembly branch instructions should be of a distinct type.

The targets for these are called **sequence symbols**.

The format of a sequence symbol is as follows.

A **sequence symbol** begins with a period (.) followed by one to seven letters or digits, the first of which must be a letter.

Unlike the symbols created by use of the **&SYSNDX** system symbol, sequence symbols do not persist into assembly time, and thus cannot generate a name conflict for the assembler.

A Sample of Conditional Assembly

Here is the DIVID macro, with conditional assembly instructions to insure that it is expanded only for parameters that are packed decimal.

```
MACRO
&LABEL  DIVID  &QUOT , &DIVIDEND , &DIVISOR
        AIF   ( T' &QUOT NE 'P' ) .NOTPACK
        AIF   ( T' &DIVIDEND NE T' &QUOT ) .NOTPACK
        AIF   ( T' &DIVISOR NE T' &QUOT ) .NOTPACK
        AGO   .DOIT
        .NOTPAK MNOTE  'ONE PARAMETER IS NOT PACKED DECIMAL'
        MEXIT
        .DOIT  ANOP
&LABEL  ZAP   &QUOT , &DIVIDEND
        CP    &DIVISOR , =P'0'  IS IT ZERO
        BNE   A&SYSNDX           NO, DIVISION IS OK
        ZAP   &QUOT , =P'0'     YES, SET QUOTIENT TO 0
        B     B&SYSNDX
A&SYSNDX DP    &QUOT , &DIVISOR
B&SYSNDX NOPR  R3              DO NOTHING
        MEND
```


Some Examples of the Conditional Assembly Divide Macro

In the following, assume that each of **X**, **Y**, and **Z** is defined by a DC statement as packed decimal, but that **A**, **B**, and **C** are defined as halfwords.

Here are some possible expansions.

```
F10DOIT  DIVID X,Y,Z
+F10DOIT  ZAP   X,Y
+        CP    Z,=P'0'           IS IT ZERO
+        BNE   A0032             NO, DIVISION IS OK
+        ZAP   X,=P'0'           YES, SET QUOTIENT TO 0
+        B     B0032
+A0032    DP    X,Z
+B0032    NOPR  R3                DO NOTHING
```

```
F25NODO  DIVID A,B,C
+ONE PARAMETER IS NOT PACKED DECIMAL
```

The Original Definition of Macro STKPOP

We now begin our redefinition of the **STKPOP** macro.

We begin with the original definition, which popped a value into a register.

***STKPOP**

```
MACRO
&L3    STKPOP &R
&L3    LH    3,STKCOUNT    GET THE STACK COUNT
        SH    3,=H'1'      SUBTRACT 1 WORD OFFSET OF TOP
        STH  3,STKCOUNT    STORE AS NEW SIZE
        SLA  3,2           BYTE OFFSET OF STACK TOP
        LA   2,THESTACK     ADDRESS OF STACK BASE
        L    &R,0(3,2)      LOAD ITEM INTO THE REGISTER.
MEND
```

Again, this macro has one symbolic parameter: **&R**. Again, a register number.

We want to expand this definition in a number of ways.

We begin by introducing the type **&TYP**.

At this point, it will become necessary to have another work register.

Mechanics of the Revised STKPOP

The new design will use register R4 to transfer the value at the top of the stack.

The new prototype will be as follows.

```
&L3      STKPOP &ARG,&TYP
```

Each type of instruction will include the following as the first statement in the “key code” – that which actually places the value into the destination.

```
      L    R4,0(3,2)    LOAD ITEM INTO REGISTER R4.
```

The second statement of the “key code” depends on the type of the destination.

```
&TYP == H      STH R4,&ARG
```

```
&TYP == F      ST  R4,&ARG
```

```
&TYP == A      ST  R4,&ARG  (SAME AS FULLWORD)
```

```
&TYP == R      LR &ARG,R4   COPY R4 INTO REGISTER
```

Again, I could define a STR macro and avoid the use of conditional assembly. For a number of reasons, I have chosen not to do so.

The Key Code as Reflected in Conditional Assembly

Again, the new prototype will be as follows.

```
&L3      STKPOP  &ARG,&TYP
```

Here is the key code section, with the conditional assembly.

The first statement is common to all types.

```
        L    R4,0(3,2)    LOAD ITEM INTO REGISTER R4.
        AIF ( `&TYPE' EQ `R' ).ISREG
        ST&TYP R4,&ARG
        AGO  .CONT
.ISREG  LR  &ARG,R4
.CONT   The next statement.
```

STKPOP: Revision 2

Here I am going to add some code to save and restore the work registers.

```
MACRO
&L3      STKPOP &ARG,&TYP
&L3      ST  R2,STKSAV2
          ST  R3,STKSAV3
          ST  R4,STKSAV4
          LH  R3,STKCOUNT      GET THE STACK COUNT
          SH  R3,=H'1'          SUBTRACT 1 WORD OFFSET OF TOP
          STH R3,STKCOUNT      STORE AS NEW SIZE
          SLA R3,2              BYTE OFFSET OF STACK TOP
          LA  R2,THESTACK       ADDRESS OF STACK BASE
          L   R4,0(3,2)         LOAD ITEM INTO REGISTER R4.
          AIF ('&TYPE' EQ 'R').ISREG
          ST&TYP R4,&ARG
          AGO .CONT
.ISREG   LR &ARG,R4
.CONT    L   R4,STKSAV4
          L   R3,STKSAV3
          L   R2,STKSAV2
MEND
```

STKPOP: The Complete Version

```
MACRO
&L3 STKPOP &ARG,&TYP
&L3 ST R3,STKSAV3
LH R3,STKCOUNT GET THE STACK COUNT
CH R3,=H'0' IS THE COUNT POSITIVE
BNH Z&SYSNDX NO, WE CANNOT POP.
SH R3,=H'1' SUBTRACT 1 WORD OFFSET OF TOP
STH R3,STKCOUNT STORE AS NEW SIZE
SLA R3,2 BYTE OFFSET OF STACK TOP
ST R2,STKSAV2 SAVE REGISTER R2
ST R4,STKSAV4 SAVE REGISTER R4
LA R2,THESTACK ADDRESS OF STACK BASE
L R4,0(3,2) LOAD ITEM INTO REGISTER R4.
AIF ('&TYPE' EQ 'R').ISREG
ST&TYP R4,&ARG
AGO .CONT
.ISREG LR &ARG,R4
.CONT L R4,STKSAV4
L R2,STKSAV2
Z&SYSNDX L R3,STKSAV3
MEND
```

Original Code for the Macro Expansions

		33 *	MACRO DEFINITIONS	
		34 *		
		35	MACRO	
		36 &L2	STKPUSH &ARG,&TYP	
		37 &L2	LH R3,STKCOUNT	
		38	SLA R3,2	
		39	LA R2,THESTACK	
		40	L&TYP R4,&ARG	
		41	ST R4,0(3,2)	
		42	LH R3,STKCOUNT	
		43	AH R3,=H'1'	
		44	STH 3,STKCOUNT	
		45	MEND	
		46 *		
		89 *	SOME MACRO INVOCATIONS	
		90 *		
		91	STKPUSH R7,R	
00004A	4830	C0C6	000CC 92+	LH R3,STKCOUNT
00004E	8B30	0002	00002 93+	SLA R3,2
000052	4120	C0CA	000D0 94+	LA R2,THESTACK
000056	1847		95+	LR R4,R7
000058	5043	2000	00000 96+	ST R4,0(3,2)
00005C	4830	C0C6	000CC 97+	LH R3,STKCOUNT
000060	4A30	C43A	00440 98+	AH R3,=H'1'
000064	4030	C0C6	000CC 99+	STH 3,STKCOUNT

000068	4830	C0C6	000CC	100	STKPUSH HHW,H
00006C	8B30	0002	000CC	101+	LH R3,STKCOUNT
000070	4120	C0CA	00002	102+	SLA R3,2
000074	4840	C1CE	000D0	103+	LA R2,THESTACK
000078	5043	2000	001D4	104+	LH R4,HHW
00007C	4830	C0C6	00000	105+	ST R4,0(3,2)
000080	4A30	C43A	000CC	106+	LH R3,STKCOUNT
000084	4030	C0C6	00440	107+	AH R3,=H'1'
			000CC	108+	STH 3,STKCOUNT
				109	STKPUSH FFW
000088	4830	C0C6	000CC	110+	LH R3,STKCOUNT
00008C	8B30	0002	00002	111+	SLA R3,2
000090	4120	C0CA	000D0	112+	LA R2,THESTACK
000094	5840	C1CA	001D0	113+	L R4,FFW
000098	5043	2000	00000	114+	ST R4,0(3,2)
00009C	4830	C0C6	000CC	115+	LH R3,STKCOUNT
0000A0	4A30	C43A	00440	116+	AH R3,=H'1'
0000A4	4030	C0C6	000CC	117+	STH 3,STKCOUNT
				118	STKPUSH FFW,A
0000A8	4830	C0E6	000EC	119+	LH R3,STKCOUNT
0000AC	8B30	0002	00002	120+	SLA R3,2
0000B0	4120	C0EA	000F0	121+	LA R2,THESTACK
0000B4	4140	C1EA	001F0	122+	LA R4,FFW
0000B8	5043	2000	00000	123+	ST R4,0(3,2)
0000BC	4830	C0E6	000EC	124+	LH R3,STKCOUNT
0000C0	4A30	C45A	00460	125+	AH R3,=H'1'
0000C4	4030	C0E6	000EC	126+	STH 3,STKCOUNT
				127 *	
				136 *****	

Revised Code for the Macros

The next few pages show the listing of the final forms of the macros, as actually coded and tested. These are followed by listings of the expanded macros.

```
002900 *
002910          MACRO
002911 &L1      STKINIT
002912 &L1      ST R3,STKSAV3
002913          SR R3,R3
002914          STH R3,STKCOUNT          CLEAR THE COUNT
002915          L  R3,STKSAV3
002920          MEND
002930 *
```

003000		MACRO	
003100	&L2	STKPUSH	&ARG,&TYP
003110	&L2	ST	R3,STKSAV3
003200		LH	R3,STKCOUNT
003210		CH	R3,STKSIZE
003220		BNL	Z&SYSNDX
003230		ST	R4,STKSAV4
003240		ST	R2,STKSAV2
003300		SLA	R3,2
003310		LA	R2,THESTACK
003320		L&TYP	R4,&ARG
003330		ST	R4,0(3,2)
003331		LH	R3,STKCOUNT
003332		AH	R3,=H'1'
003333		STH	3,STKCOUNT
003334		L	R4,STKSAV4
003335		L	R2,STKSAV2
003336	Z&SYSNDX	L	R3,STKSAV3
003337		MEND	
003338	*		
003339	*		

SAVE REGISTER R3
GET THE CURRENT SIZE
IS THE STACK FULL?
YES, DO NOT PUSH
OK, SAVE R2 AND R4

MULTIPLY BY FOUR
ADDRESS OF STACK START
LOAD R4 WITH VALUE
STORE INTO THE STACK

```

003340          MACRO
003341 &L3        STKPOP &ARG,&TYP
003342 &L3        ST      R3,STKSAV3
003343          LH      R3,STKCOUNT          GET THE STACK COUNT
003344          CH      R3,=H'0'             IS THE COUNT POSITIVE?
003345          BNH     Z&SYSNDX             NO, WE CANNOT POP
003346          SH      R3,=H'1'             SUBTRACT 1 WORD OFFSET
003347          STH     R3,STKCOUNT         STORE THE NEW SIZE
003348          SLA     R3,2                 BYTE OFFSET OF STACK TOP
003349          ST      R2,STKSAV2           SAVE REGISTER R2
003350          ST      R4,STKSAV4           SAVE REGISTER R4
003351          LA      R2,THESTACK         ADDRESS OF STACK BASE
003352          L       R4,0(3,2)           LOAD ITEM INTO R4
003353          AIF    ('&TYP' EQ 'R').ISREG
003354          ST&TYP R4,&ARG
003355          AGO    .CONT
003356 .ISREG    LR &ARG,R4
003357 .CONT     L     R4,STKSAV4
003358          L     R2,STKSAV2
003359 Z&SYSNDX L     R3,STKSAV3
003360          MEND
003361 *

```

Revised Code for the Macro Expansions

			128 *	SOME MACRO INVOCATIONS
			129 *	
			130	STKINIT
00004A	5030	C22E	00234 131+	ST R3,STKSAV3
00004E	1B33		132+	SR R3,R3
000050	4030	C226	0022C 133+	STH R3,STKCOUNT
000054	5830	C22E	00234 134+	L R3,STKSAV3
			135 *	
			136	STKPUSH R7,R
000058	5030	C22E	00234 137+	ST R3,STKSAV3
00005C	4830	C226	0022C 138+	LH R3,STKCOUNT
000060	4930	C228	0022E 139+	CH R3,STKSIZE
000064	47B0	C08C	00092 140+	BNL Z0010
000068	5040	C232	00238 141+	ST R4,STKSAV4
00006C	5020	C22A	00230 142+	ST R2,STKSAV2
000070	8B30	0002	00002 143+	SLA R3,2
000074	4120	C236	0023C 144+	LA R2,THESTACK
000078	1847		145+	LR R4,R7
00007A	5043	2000	00000 146+	ST R4,0(3,2)
00007E	4830	C226	0022C 147+	LH R3,STKCOUNT
000082	4A30	C5A2	005A8 148+	AH R3,=H'1'
000086	4030	C226	0022C 149+	STH 3,STKCOUNT
00008A	5840	C232	00238 150+	L R4,STKSAV4
00008E	5820	C22A	00230 151+	L R2,STKSAV2
000092	5830	C22E	00234 152+Z0010	L R3,STKSAV3

000096	5030	C22E	00234	153	STKPUSH	HHW,H
00009A	4830	C226	00234	154+	ST	R3,STKSAV3
00009E	4930	C228	0022C	155+	LH	R3,STKCOUNT
0000A2	47B0	C0CC	0022E	156+	CH	R3,STKSIZE
0000A6	5040	C232	000D2	157+	BNL	Z0011
0000AA	5020	C22A	00238	158+	ST	R4,STKSAV4
0000AE	8B30	0002	00230	159+	ST	R2,STKSAV2
0000B2	4120	C236	00002	160+	SLA	R3,2
0000B6	4840	C33A	0023C	161+	LA	R2,THESTACK
0000BA	5043	2000	00340	162+	LH	R4,HHW
0000BE	4830	C226	00000	163+	ST	R4,0(3,2)
0000C2	4A30	C5A2	0022C	164+	LH	R3,STKCOUNT
0000C6	4030	C226	005A8	165+	AH	R3,=H'1'
0000CA	5840	C232	0022C	166+	STH	3,STKCOUNT
0000CE	5820	C22A	00238	167+	L	R4,STKSAV4
0000D2	5830	C22E	00230	168+	L	R2,STKSAV2
			00234	169+	L	R3,STKSAV3

			170		STKPUSH FFW
0000D6	5030	C22E	00234	171+	ST R3,STKSAV3
0000DA	4830	C226	0022C	172+	LH R3,STKCOUNT
0000DE	4930	C228	0022E	173+	CH R3,STKSIZE
0000E2	47B0	C10C	00112	174+	BNL Z0012
0000E6	5040	C232	00238	175+	ST R4,STKSAV4
0000EA	5020	C22A	00230	176+	ST R2,STKSAV2
0000EE	8B30	0002	00002	177+	SLA R3,2
0000F2	4120	C236	0023C	178+	LA R2,THESTACK
0000F6	5840	C336	0033C	179+	L R4,FFW
0000FA	5043	2000	00000	180+	ST R4,0(3,2)
0000FE	4830	C226	0022C	181+	LH R3,STKCOUNT
000102	4A30	C5A2	005A8	182+	AH R3,=H'1'
000106	4030	C226	0022C	183+	STH 3,STKCOUNT
00010A	5840	C232	00238	184+	L R4,STKSAV4
00010E	5820	C22A	00230	185+	L R2,STKSAV2
000112	5830	C22E	00234	186+ Z0012	L R3,STKSAV3

000116	5030	C22E	00234	187	STKPUSH FFW,A
00011A	4830	C226	00234	188+	ST R3,STKSAV3
00011E	4930	C228	0022C	189+	LH R3,STKCOUNT
000122	47B0	C14C	0022E	190+	CH R3,STKSIZE
000126	5040	C232	00152	191+	BNL Z0013
00012A	5020	C22A	00238	192+	ST R4,STKSAV4
00012E	8B30	0002	00230	193+	ST R2,STKSAV2
000132	4120	C236	00002	194+	SLA R3,2
000136	4140	C336	0023C	195+	LA R2,THESTACK
00013A	5043	2000	0033C	196+	LA R4,FFW
00013E	4830	C226	00000	197+	ST R4,0(3,2)
000142	4A30	C5A2	0022C	198+	LH R3,STKCOUNT
000146	4030	C226	005A8	199+	AH R3,=H'1'
00014A	5840	C232	0022C	200+	STH 3,STKCOUNT
00014E	5820	C22A	00238	201+	L R4,STKSAV4
000152	5830	C22E	00230	202+	L R2,STKSAV2
			00234	203+Z0013	L R3,STKSAV3

			204 *		
			205		STKPOP R8,R
000156	5030	C22E	00234	206+	ST R3,STKSAV3
00015A	4830	C226	0022C	207+	LH R3,STKCOUNT
00015E	4930	C5A4	005AA	208+	CH R3,=H'0'
000162	47D0	C186	0018C	209+	BNH Z0014
000166	4B30	C5A2	005A8	210+	SH R3,=H'1'
00016A	4030	C226	0022C	211+	STH R3,STKCOUNT
00016E	8B30	0002	00002	212+	SLA R3,2
000172	5020	C22A	00230	213+	ST R2,STKSAV2
000176	5040	C232	00238	214+	ST R4,STKSAV4
00017A	4120	C236	0023C	215+	LA R2,THESTACK
00017E	5843	2000	00000	216+	L R4,0(3,2)
000182	1884			217+	LR R8,R4
000184	5840	C232	00238	218+	L R4,STKSAV4
000188	5820	C22A	00230	219+	L R2,STKSAV2
00018C	5830	C22E	00234	220+ Z0014	L R3,STKSAV3

000190	5030	C22E	00234	221	STKPOP	FFW
000194	4830	C226	00234	222+	ST	R3,STKSAV3
000198	4930	C5A4	0022C	223+	LH	R3,STKCOUNT
00019C	47D0	C1C2	005AA	224+	CH	R3,=H'0'
0001A0	4B30	C5A2	001C8	225+	BNH	Z0015
0001A4	4030	C226	005A8	226+	SH	R3,=H'1'
0001A8	8B30	0002	0022C	227+	STH	R3,STKCOUNT
0001AC	5020	C22A	00002	228+	SLA	R3,2
0001B0	5040	C232	00230	229+	ST	R2,STKSAV2
0001B4	4120	C236	00238	230+	ST	R4,STKSAV4
0001B8	5843	2000	0023C	231+	LA	R2,THESTACK
0001BC	5040	C336	00000	232+	L	R4,0(3,2)
0001C0	5840	C232	0033C	233+	ST	R4,FFW
0001C4	5820	C22A	00238	234+	L	R4,STKSAV4
0001C8	5830	C22E	00230	235+	L	R2,STKSAV2
			00234	236+Z0015	L	R3,STKSAV3

0001CC	5030	C22E	00234	237	STKPOP	HHW,H
0001D0	4830	C226	00234	238+	ST	R3,STKSAV3
0001D4	4930	C5A4	0022C	239+	LH	R3,STKCOUNT
0001D8	47D0	C1FE	005AA	240+	CH	R3,=H'0'
0001DC	4B30	C5A2	00204	241+	BNH	Z0016
0001E0	4030	C226	005A8	242+	SH	R3,=H'1'
0001E4	8B30	0002	0022C	243+	STH	R3,STKCOUNT
0001E8	5020	C22A	00002	244+	SLA	R3,2
0001EC	5040	C232	00230	245+	ST	R2,STKSAV2
0001F0	4120	C236	00238	246+	ST	R4,STKSAV4
0001F4	5843	2000	0023C	247+	LA	R2,THESTACK
0001F8	4040	C33A	00000	248+	L	R4,0(3,2)
0001FC	5840	C232	00340	249+	STH	R4,HHW
000200	5820	C22A	00238	250+	L	R4,STKSAV4
000204	5830	C22E	00230	251+	L	R2,STKSAV2
			00234	252+Z0016	L	R3,STKSAV3
				253 *		

Revised Code for the Macro STKINIT

Here is an expansion of the newer definition of STKINIT,
which allows the stack size to be specified.

00004A 5030 C05E	00064	138	STKINIT 128
00004E 1B33		139+	ST R3,STKSAV3
000050 4030 C056	0005C	140+	SR R3,R3
000054 5830 C05E		141+	STH R3,STKCOUNT
000058 47F0 C266	00064	142+	L R3,STKSAV3
00005C 0000	0026C	143+	B L0009
00005E 0080		144+STKCOUNT	DC H'0'
000060 00000000		145+STKSIZE	DC H'128'
000064 00000000		146+STKSAV2	DC F'0'
000068 00000000		147+STKSAV3	DC F'0'
00006C 0000000000000000		148+STKSAV4	DC F'0'
00026C 8B30 0000	00000	149+THESTACK	DC 128F'0'
		150+L0009	SLA R3,0